



Version 4.0 was released in July 2013

Now available in most production compilers

What's new in 4.0



- User defined reductions
- Construct cancellation
- Portable SIMD directives
- Extensions to tasking
- Thread affinity
- Accelerator offload support

User defined reductions



- As of 3.1 cannot do reductions on objects or structures.
- UDR extensions in 4.0 add support for this.

- Use declare reduction directive to define new reduction operators
- New operators can then be used in reduction clause.

```
#pragma omp declare reduction (reduction-identifier :
typename-list : combiner) [identity(identity-expr)]
```



- reduction-identifier gives a name to the operator
 - Can be overloaded for different types
 - Can be redefined in inner scopes
- typename-list is a list of types to which it applies
- combiner expression specifies how to combine values
- identity can specify the identity value of the operator
 Can be an expression or a brace initializer

```
#pragma omp declare reduction (merge : std::vector<int>
: omp_out.insert(omp_out.end(), omp_in.begin(), omp_in.end()))
```

- Private copies created for a reduction are initialized to the identity that was specified for the operator and type
 - Default identity defined if identity clause not present
- Compiler uses combiner to combine private copies
- omp_out refers to private copy that holds combined values
- omp_in refers to the other private copy
- Can now use merge as a reduction operator.

Construct cancellation



- Clean way to signal early termination of an OpenMP construct.
 - one thread signals
 - other threads jump to the end of the construct

```
!$omp cancel construct [if (expr)]
```

where *construct* is **parallel**, **sections**, **do** or **taskgroup** cancels the construct

!\$omp cancellation point construct checks for cancellation (also happens implicitly at cancel directive, barriers etc.)



```
!$omp parallel do private(eureka)
do i=1,n
    eureka = testing(i,...)
!$omp cancel parallel if(eureka)
end do
```

- First thread for which eureka is true will cancel the parallel region and exit.
- Other threads exit next time they hit the cancel directive

Portable SIMD directives



- Many compilers support SIMD directives to aid vectorisation of loops.
 - compiler can struggle to generate SIMD code without these
- OpenMP 4.0 provides a standardised set
- Use simd directive to indicate a loop should be SIMDized

```
#pragma omp simd [clauses]
```

- Executes iterations of following loop in SIMD chunks
- Loop is not divided across threads
- SIMD chunk is set of iterations executed concurrently by SIMD lanes



- Clauses control data environment, how loop is partitioned
- safelen (length) limits the number of iterations in a SIMD chunk.
- linear lists variables with a linear relationship to the iteration space
- aligned specifies byte alignments of a list of variables
- private, lastprivate, reduction and collapse have usual meanings.
- Also declare simd directive to generate SIMDised versions of functions.
- Can be combined with loop constructs (parallelise and SIMDise)

Extensions to tasking



- taskgroup directive provide allow task to wait for all descendant tasks to complete
- Compare taskwait, which only waits for children

```
#pragma omp taskgroup
{
    create_a_group_of_tasks(could_create_nested_tasks);
} // all created tasks complete by here
```

Task dependencies



depend clause on task construct

!\$omp task depend(type: list)

where type is in, out or inout and list is a list of variables.

- list may contain subarrays: OpenMP 4.0 includes a syntax for C/C++
- in: the generated task will be a dependent task of all previously generated sibling tasks that reference at least one of the list items in an out or inout clause.
- out or inout: the generated task will be a dependent task of all previously generated sibling tasks that reference at least one of the list items in in, out or inout clause.



```
#pragma omp task depend (out:a)
    { ... }

#pragma omp task depend (out:b)
    { ... }

#pragma omp task depend (in:a,b)
    { ... }
```

- The first two tasks can execute in parallel
- The third task cannot start until both the first two are complete

Thread affinity



- Since many systems are now NUMA and SMT, placement of threads on the hardware can have a big effect on performance.
- Up until now, control of this in OpenMP is very limited.
- Some compilers have their own extensions.
- OpenMP 4.0 gives much more control

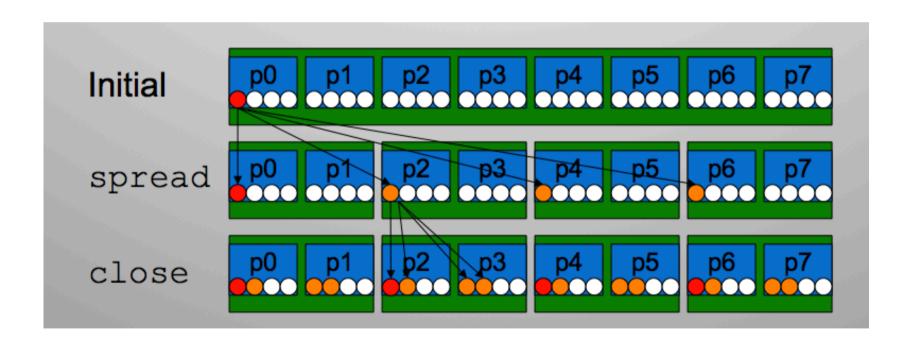
Affinity environment



- Increased choices for OMP_PROC_BIND
- Can still specify true or false
- Can now provide a list (possible item values: master,
 close or spread) to specify how to bind parallel regions at different nesting levels.
- Added OMP_PLACES environment variable
- Can specify abstract names including threads, cores and sockets
- Can specify an explicit ordered list of places
- Place numbering is implementation defined



export OMP_PLACES=threads
export OMP_PROC_BIND="spread,close"



Accelerator support



- Similar to, but not the same as, OpenACC directives.
- Support for more than just loops
- Less reliance on compiler to parallelise and map code to threads
- Not GPU specific
- Fully integrated into OpenMP



- Host-centric model with one host device and multiple target devices of the same type.
- device: a logical execution engine with local storage.
- device data environment: a data environment associated with a target data or target region.
- target constructs control how data and code is offloaded to a device.
- Data is mapped from a host data environment to a device data environment.

- Code inside target region is executed on the device.
- Executes sequentially by default.
- Can include other OpenMP directives to run in parallel
- Clauses to control data movement.

```
#pragma omp target map(to:B,C), map(tofrom:sum)
#pragma omp parallel for reduction(+:sum)
for (int i=0; i<N; i++) {
    sum += B[i] + C[i];
}</pre>
```



- target data construct just moves data and does not execute code (c.f. #pragma acc data in OpenACC.
- target update construct updates data during a target data region.
- declare target compiles a version of function/subroutine that can be called on the device.
- Target regions are blocking: the encountering thread waits for them to complete.
 - Asynchronous behaviour can be achieved by using target regions inside tasks (with dependencies if required).

What about GPUs?



- Executing a target region on a GPU can only use one multiprocessor
 - synchronisation required for OpenMP not possible between multiprocessors
 - not much use!
- teams construct creates multiple master threads which can execute in parallel, spawn parallel regions, but not synchronise or communicate with each other.
- distribute construct spreads the iterations of a parallel loop across teams.
 - Only schedule option is static (with optional chunksize).

```
#pragma omp target teams distribute parallel for\
map(to:B,C), map(tofrom:sum) reduction(+:sum)

for (int i=0; i<N; i++) {
    sum += B[i] + C[i];
}</pre>
```

• Distributes iterations across multiprocessors *and* across threads within each multiprocessor.